

Recursive Least-Squares Algorithm

Let

$$z(t) = \sum_{i=1}^N w_i r_i(t) \quad (1)$$

with the task of making $z(t) \approx f(t)$. The least-square solution is

$$w_i = \sum_{j=1}^N P_{ij} D_j \quad (2)$$

where

$$P_{ij} = (C^{-1})_{ij} \quad \text{and} \quad D_j = \langle r_j f \rangle \quad (3)$$

with

$$C_{ij} = \langle r_i r_j \rangle. \quad (4)$$

Doing this recursively, we write

$$\begin{aligned} w_i(t) &= w_i(t-1) + \Delta w_i(t) \\ C_{ij}(t) &= C_{ij}(t-1) + r_i(t)r_j(t) \\ P_{ij}(t) &= P_{ij}(t-1) + \Delta P_{ij}(t) \\ D_j(t) &= D_j(t-1) + r_j(t)f(t). \end{aligned} \quad (5)$$

In the following, r_i always means $r_i(t)$ and f means $f(t)$. Now,

$$\Delta P_{ij}(t) = -\frac{(\sum_k P_{ik}(t-1)r_k)(\sum_l r_l P_{lj}(t-1))}{1 + \sum_{k,l} r_k P_{kl}(t-1)r_l}. \quad (6)$$

A useful identity arising from this results is

$$\sum_j P_{ij}(t)r_j = \frac{\sum_j P_{ij}(t-1)r_j}{1 + \sum_{k,l} r_k P_{kl}(t-1)r_l}. \quad (7)$$

We can write

$$\Delta w_i(t) = \sum_j (\Delta P_{ij}(t)D_j(t-1) + P_{ij}(t)r_j f) \quad (8)$$

and

$$\sum_j \Delta P_{ij}(t)D_j(t-1) = -\frac{(\sum_k P_{ik}(t-1)r_k)(\sum_{j,l} r_l P_{lj}(t-1)D_j(t-1))}{1 + \sum_{k,l} r_k P_{kl}(t-1)r_l} \quad (9)$$

but

$$\sum_{j,l} r_l P_{lj}(t-1)D_j(t-1) = \sum_l w_l(t-1)r_l. \quad (10)$$

so

$$\sum_j \Delta P_{ij}(t) D_j(t-1) = - \frac{(\sum_k P_{ik}(t-1)r_k) (\sum_l w_l(t-1)r_l)}{1 + \sum_{k,l} r_k P_{kl}(t-1)r_l} \quad (11)$$

Using equation 7, this gives

$$\sum_j \Delta P_{ij}(t) D_j(t-1) = - \sum_j P_{ij}(t) r_j \sum_k w_k(t-1) r_k . \quad (12)$$

Putting equations 8 and 12 together.

$$\Delta w_i(t) = \sum_j P_{ij}(t) r_j e(t) \quad (13)$$

where

$$e(t) = f(t) - \sum_k w_k(t-1) r_k(t) \quad (14)$$

and with equation 6 giving $P_{ij}(t)$.